

STABILIZED COLUMN GENERATION

Jacques Desrosiers

HEC & GERAD

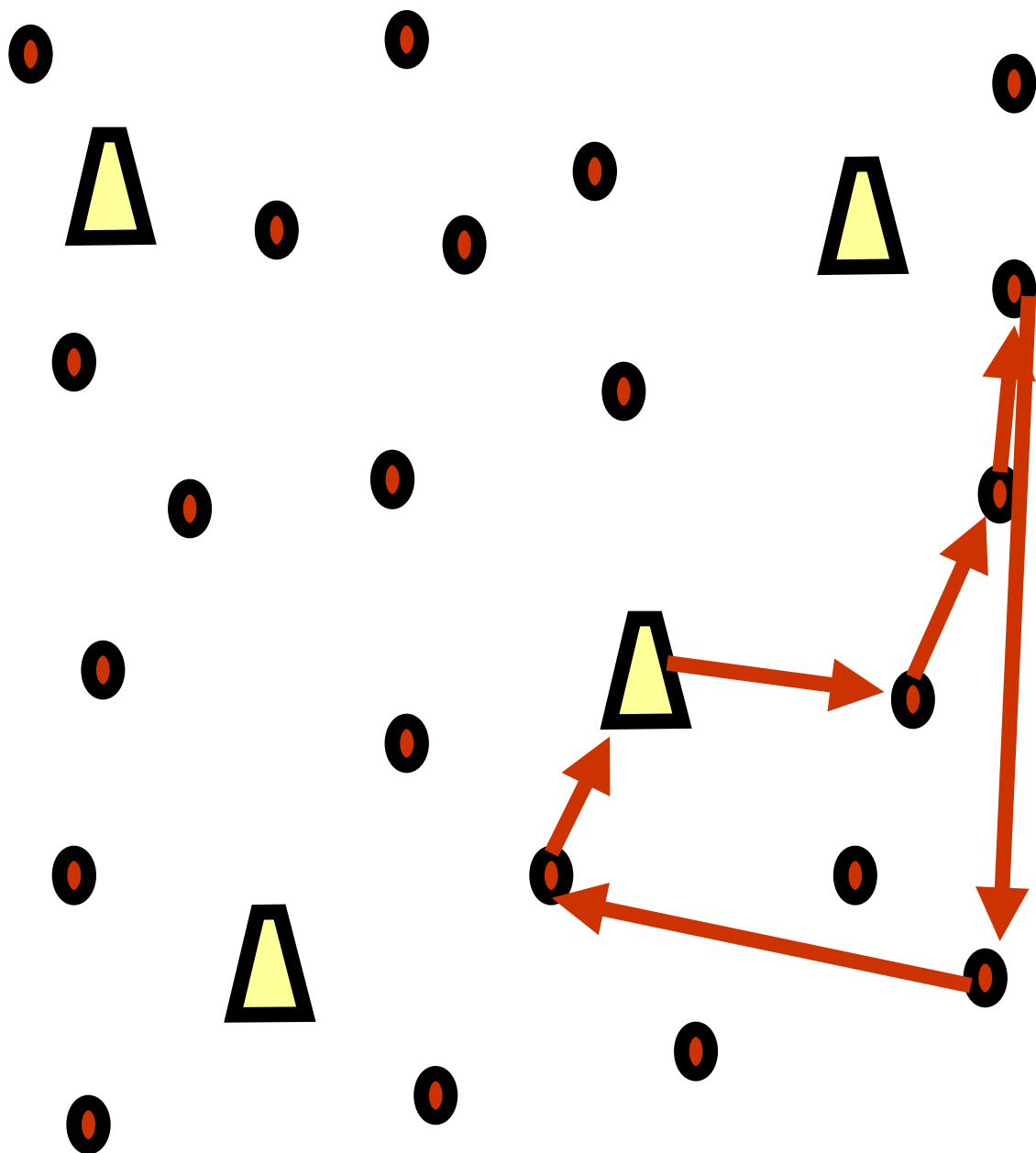
Hatem Ben Amor

François Soumis

Daniel Villeneuve

GERAD

Multiple Depot Vehicle Scheduling Problem



MDVSP

K : set of depots

N : set of scheduled tasks (no time windows)

t_i : time at which service must start for task $i \in N$

s_i : service duration for task $i \in N$

d_{ij} : time to go from node i to node j , for $i, j \in N \cup K$

A : set of arcs; $(i, j) \in A$ if $t_i + s_i + d_{ij} \leq t_j$, for $i, j \in N \cup K$

Ω^k : set of paths from and to depot $k \in K$

c_p^k : cost of path $p \in \Omega^k$

$a_{ip}^k = 1$ if path $p \in \Omega^k$ covers task $i \in N$; 0 otherwise

n^k : available number of vehicles at depot $k \in K$

$$\min \sum_{k \in K} \sum_{p \in \Omega^k} c_p^k x_p^k$$

$$\sum_{k \in K} \sum_{p \in \Omega^k} a_{ip}^k x_p^k = 1, \quad i \in N$$

$$\sum_{p \in \Omega^k} x_p^k \leq n^k, \quad k \in K$$

$$x_p^k \text{ binary}, \quad k \in K, p \in \Omega^k$$

Motivation

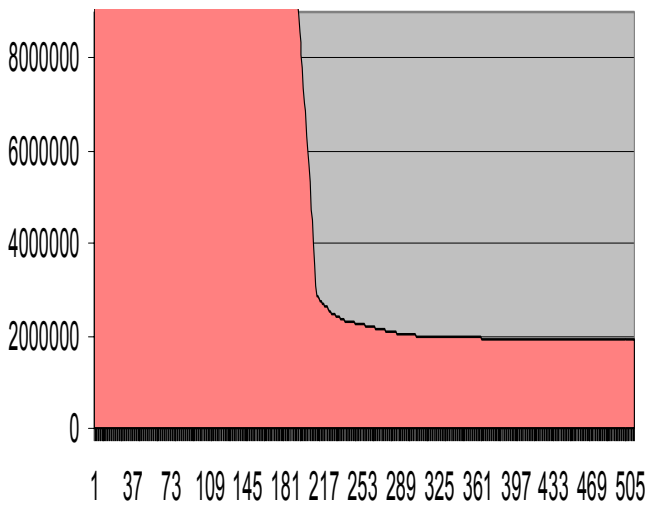
<i>MDVSP</i>	cpu	cpu	cpu	# CG	# SP	# MP
<i>R 800 (4)</i>	tot	mp	sp	iter	cols	itr
<i>standard</i>	4178.4	3149.2	1029.2	509	37579	926161

Tailing off effect of the objective function

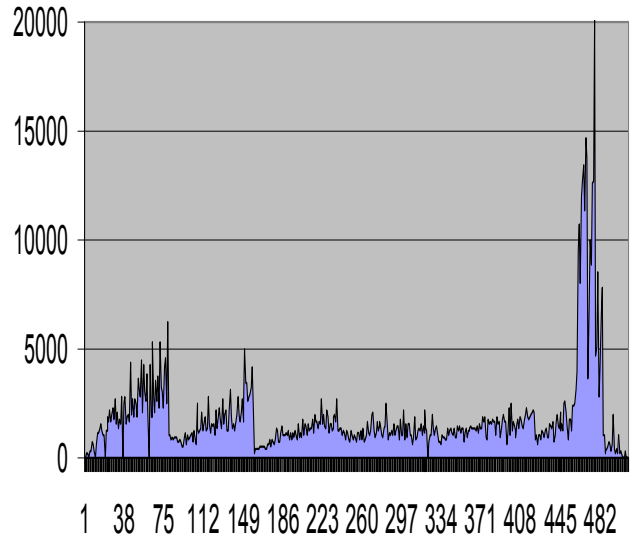


1 33 65 97 129 161 193 225 257 289 321 353 385 417 449 481

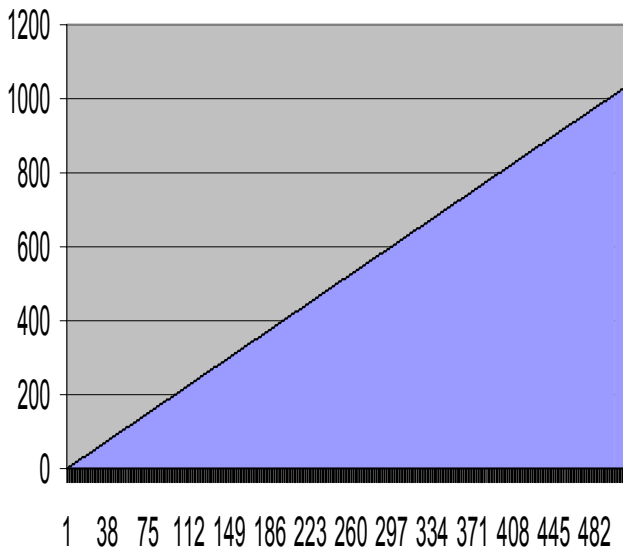
MP Objective



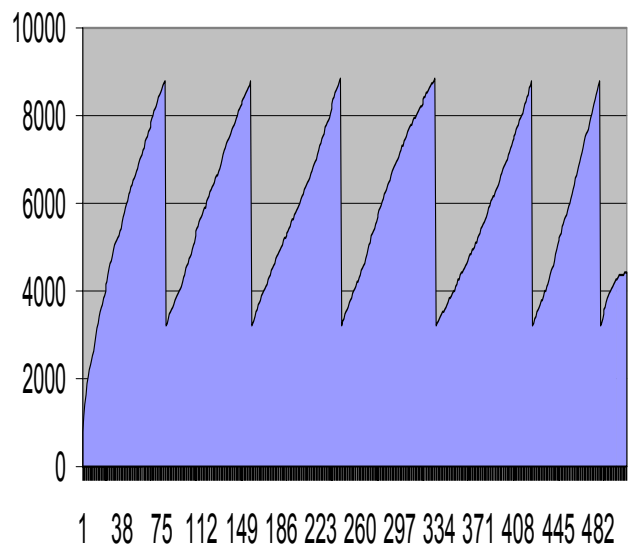
MP Iterations



SP CPU Time



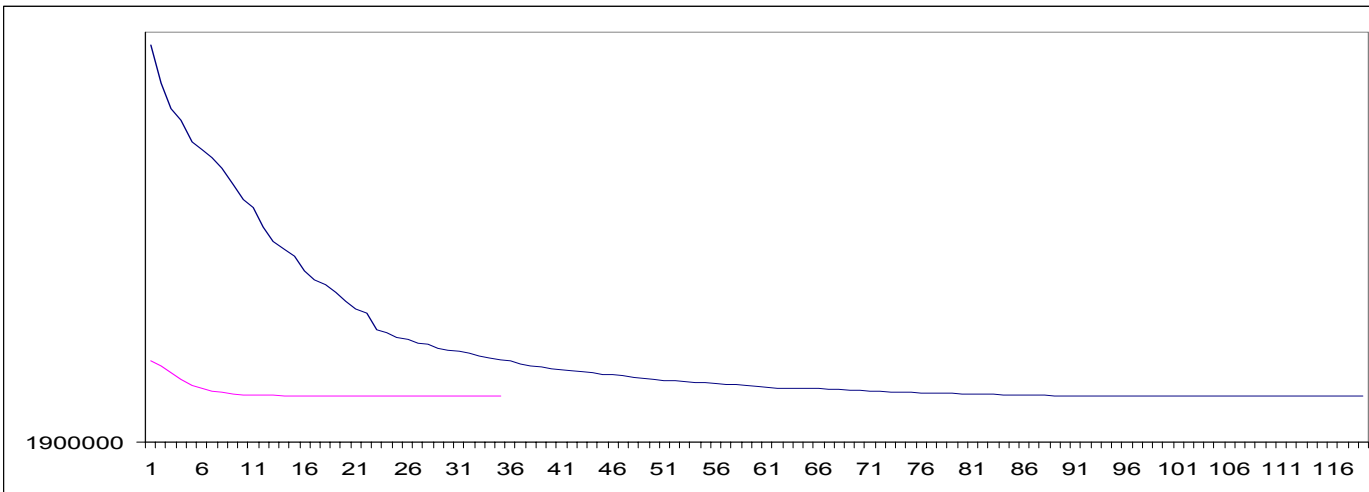
MP Columns



Using Optimal Dual Values

Problem			cpu	# CG	# SP	# MP
R 800 (4)	Opt sol	Init sol	tot	iter	cols	itr
standard	1915589.5	800000000	4178.4	509	37579	926161
delta boxes						
100		2035590.5	835.5	119	9368	279155
10		1927590.5	117.9	35	2789	40599
1		1916790.5	52.0	20	1430	8744
0.1		1915710.5	47.5	19	1333	8630
0.01	1915589.1	1915602.5	37.3	17	1145	6288

delta boxes	Final gap(%)	Initial gap (%)	(results in % vs standard set)			
100		6.26	20.0	23.4	24.9	30.1
10		0.63	2.8	6.9	7.4	4.4
1		0.063	1.2	3.9	3.8	0.9
0.1		0.0063	1.1	3.7	3.5	0.9
0.01	-0.00002088	0.0007	0.9	3.3	3.0	0.7



Optimal Dual Values

- **Useful in the context of Lagrangian Relaxation to recover primal feasibility**
- **Useful to perform crossover from an interior point solution to an extreme point solution**

Crossover:

CPLEX vs STABILIZATION

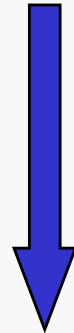
(cpu times in seconds)

	Problem 1	Problem 2	Problem 3	Problem 4	Problem 5
Constraints	12,354	12,313	13,344	13,453	13,269
Variables	126,329	129,349	151,665	156,841	148,040
CPLEX Primopt	> 10000	> 10000	> 10000	> 10000	> 10000
CPLEX Baropt	1,056	1,011	1,342	1,382	1,171
CROSSOVER					
CPLEX Promopt	13	3,427	> 10000	97	6,659
CPLEX Dualopt	50	697	1,876	686	778
Stabilazation 10 ⁻¹	189	291	671	631	417
Stabilazation 10 ⁻²	100	121	554	557	380
Stabilazation 10 ⁻³	91	102	407	396	353
Stabilazation 10 ⁻⁴	87	92	350	347	285

LP Column Generation

MASTER PROBLEM

Columns



Dual Multipliers

COLUMN GENERATORS

(Shortest Path Problems on Acyclic Graphs)

Optimality Conditions:

primal feasibility

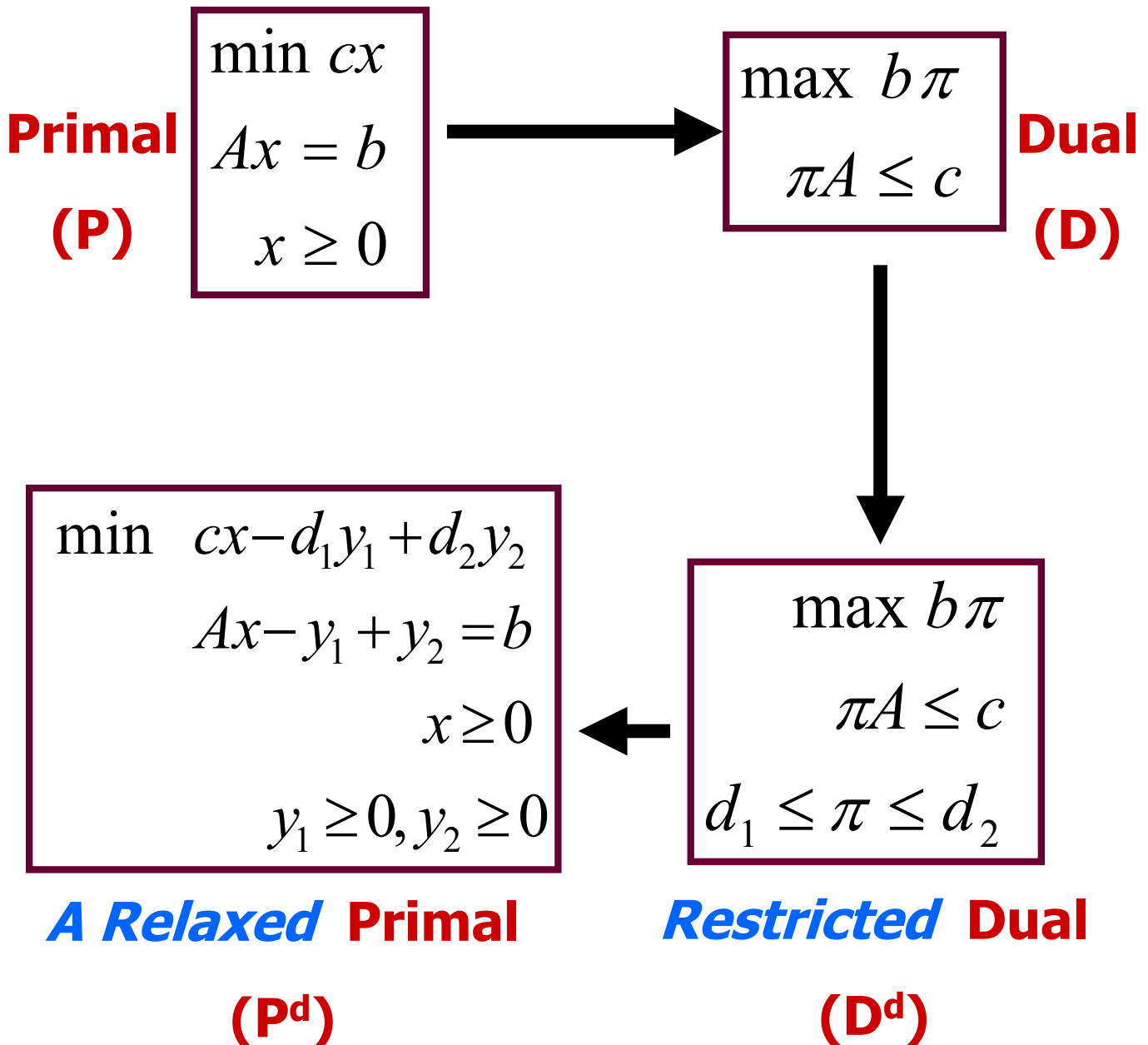
dual feasibility

complementary slackness

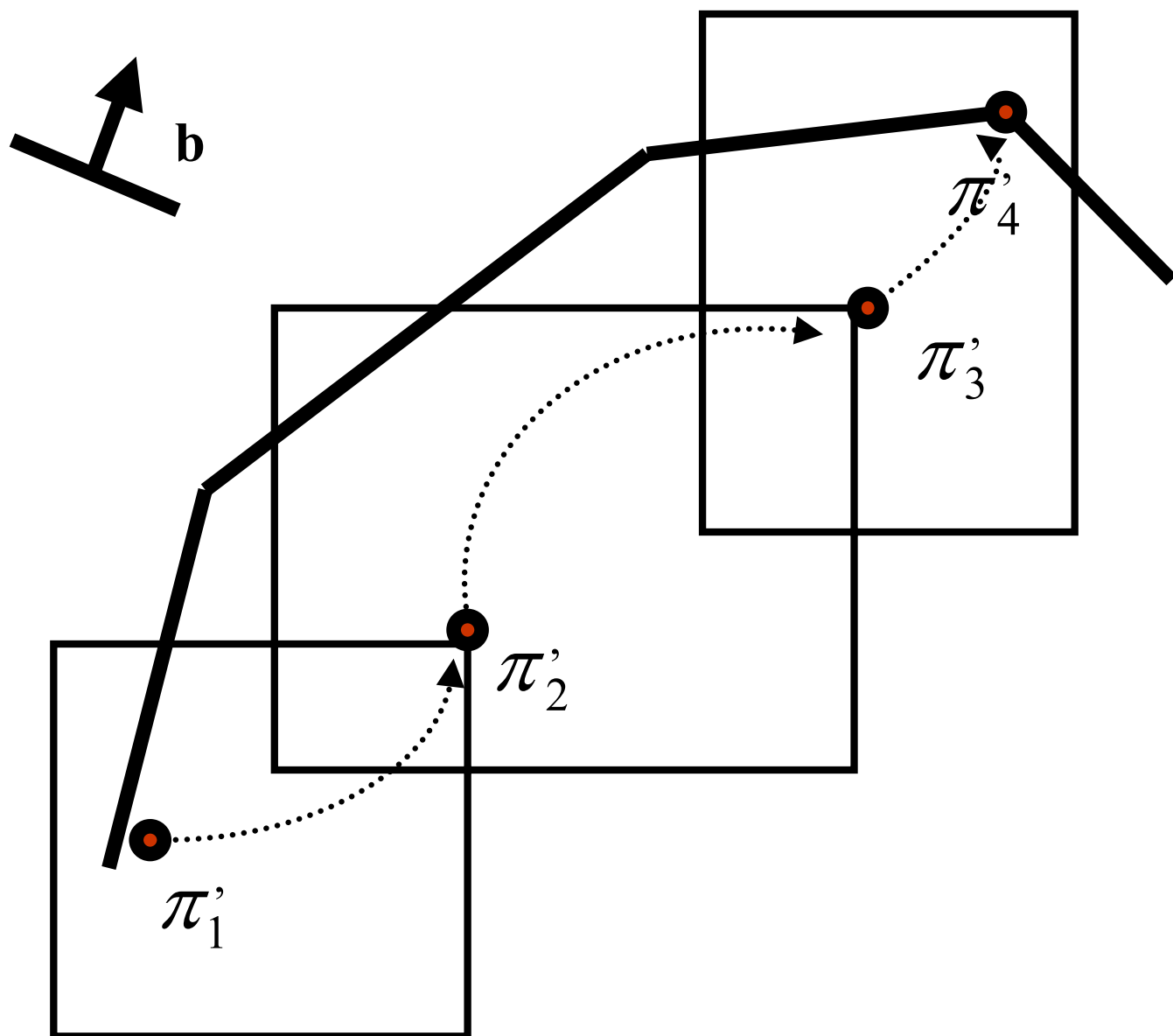
Some References

1960 Gillmore & Gomory	Cutting Stock Problem
1989 Agarwal, Mathur & Salkin	VRP
1963 Marquardt	Trust region
1975 Marsten	Box Step
2000 Kallehauge & Madsen	VRPTW
1992 Vial & Goffin	ACCPM
1999 du Merle et al.	Stabilized CG
2000 Valério de Carvalho	Cutting Stock Problem

Impact of Dual Bounds on the Primal Formulation



Trust Regions / Box Step



Dual Cuts

for the Cutting Stock problem

Valério de Carvalho (2000)

*“Using extra **dual cuts** to accelerate column generation”*

Small items ($i=1,\dots,m$) are ranked :

$$l_1 > l_2 > l_3 > \dots \implies \pi_1 \geq \pi_2 \geq \pi_3 \geq \dots$$

Additionally:

$$l_i \geq l_j + l_k \implies \pi_i \geq \pi_j + \pi_k$$

Using a priori at most $2m$ dual constraints
(or primal columns)

...

Dual Cuts / Primal Columns

Valério de Carvalho (2000)

Application to the Cutting Stock Problem

Generated cutting patterns			a priori columns				
			1			1	
			-1	1			1
				-1	1		
					-1		
					...		
						-1	
						-1	-1
							-1
							...

reduces the CPU time by 40%.

Triplets (501 items)

each roll of length L

is cut into exactly three orders without any waste

Standard Column Generation	124.2 iterations
----------------------------	------------------

$\pi_1 \geq \pi_2 \geq \pi_3 \geq \dots$	113.3 iterations
--	------------------

$\pi_i = \ell_i / L, \quad i = 1, \dots, m$	12.2 iterations
---	-----------------

Both strategies	11.2 iterations
-----------------	-----------------

Average over 10 problems

Degeneracy & Perturbation

Primal
(P)

$$\begin{aligned} \min \quad & cx \\ \text{subject to} \quad & Ax = b \\ & x \geq 0 \end{aligned}$$

$$\begin{aligned} \min \quad & cx \\ \text{subject to} \quad & Ax - y_1 + y_2 = b \\ & x \geq 0 \\ & 0 \leq y_1 \leq \varepsilon_1, 0 \leq y_2 \leq \varepsilon_2 \end{aligned}$$

***Perturbed* Primal**

(P_e)

A relaxed primal

Stabilized Column Generation

$$\min cx$$

$$Ax = b$$

$$x \geq 0$$

$$\max b\pi$$

$$\pi A \leq c$$

$$\min cx$$

$$Ax - y_1 + y_2 = b$$

$$x \geq 0$$

$$0 \leq y_1 \leq \varepsilon_1, 0 \leq y_2 \leq \varepsilon_2$$

$$\max b\pi$$

$$\pi A \leq c$$

$$d_1 \leq \pi \leq d_2$$

Perturbed Primal

Restricted Dual

$$\min \quad cx - d_1 y_1 + d_2 y_2$$

$$Ax - y_1 + y_2 = b$$

$$x \geq 0$$

$$0 \leq y_1 \leq \varepsilon_1, \quad 0 \leq y_2 \leq \varepsilon_2$$

Stabilized Primal (P_ε^d)

Stabilized Primal & Dual

$$\min \quad cx - d_1 y_1 + d_2 y_2$$

$$Ax - y_1 + y_2 = b \quad \pi$$

$$y_1 \leq \varepsilon_1 \quad -\omega_1 \leq 0$$

$$y_2 \leq \varepsilon_2 \quad -\omega_2 \leq 0$$

$$x \geq 0, y_1 \geq 0, y_2 \geq 0$$

Stabilized Primal (P_ε^d)

$$\max \quad b\pi - \omega_1 \varepsilon_1 - \omega_2 \varepsilon_2$$

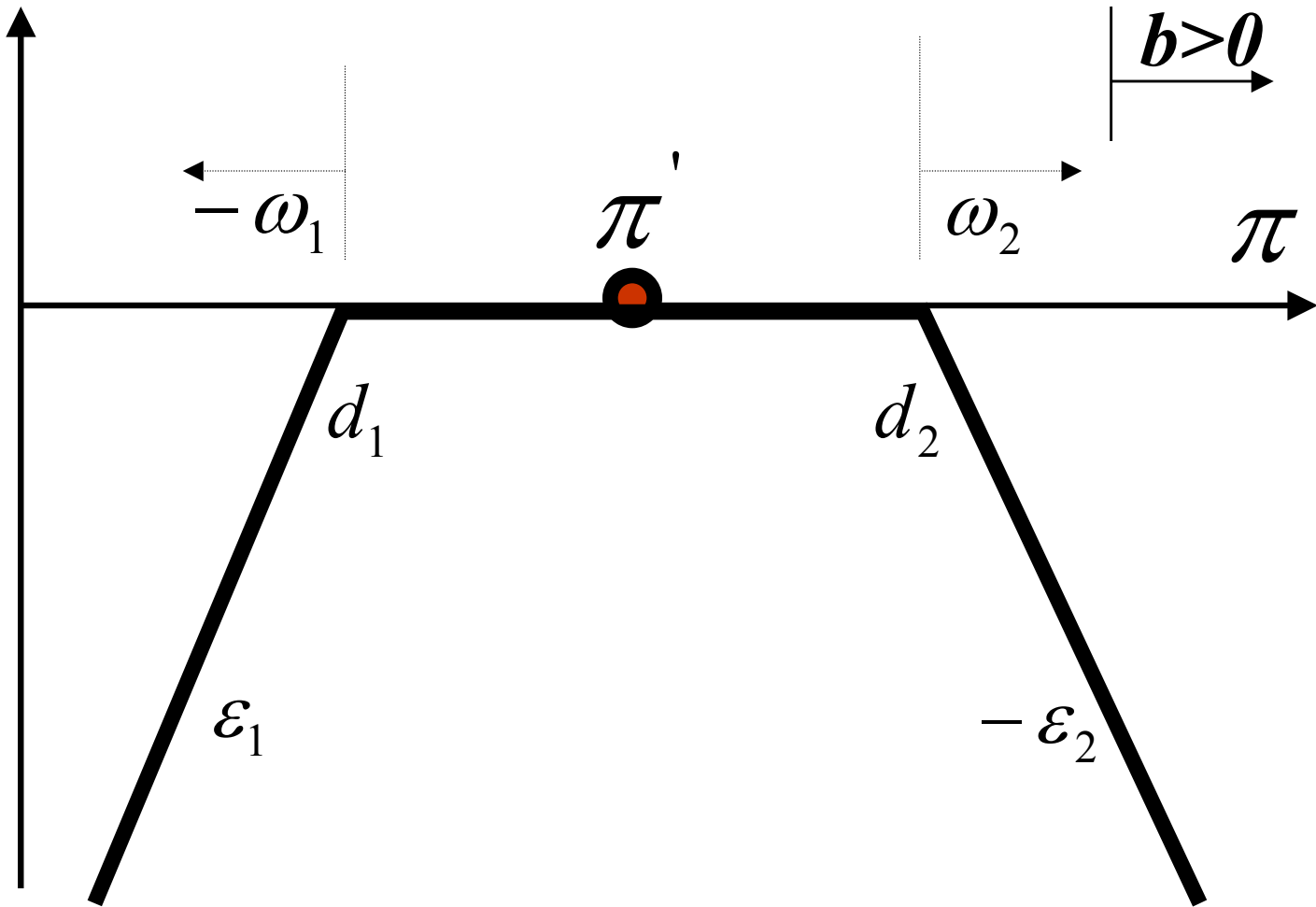
$$\pi A \leq c$$

$$d_1 - \omega_1 \leq \pi \leq d_2 + \omega_2$$

$$\omega_1 \geq 0, \omega_2 \geq 0$$

Stabilized Dual (D_ε^d)

Dual Space



$$d_1 - \omega_1 \leq \pi \leq d_2 + \omega_2$$
$$\omega_1 \geq 0, \omega_2 \geq 0$$

Propositions

Within a column generation scheme,
problem P_ε^d is optimal when
 $\bar{c}_p \geq 0$, for all columns $p \in \Omega$

Since P_ε^d is a relaxation of P
 $cx - d_1y_1 + d_2y_2 = b\pi - \varepsilon_1\omega_1 - \varepsilon_2\omega_2$
is a lower bound on $v(P)$.

$$b\pi$$

is a better lower bound on $v(P)$.

Propositions

Problem $P \equiv P_\varepsilon^d$ if
 $\varepsilon = 0$, or $\exists \pi^* \mapsto d_1 < \pi^* < d_2$

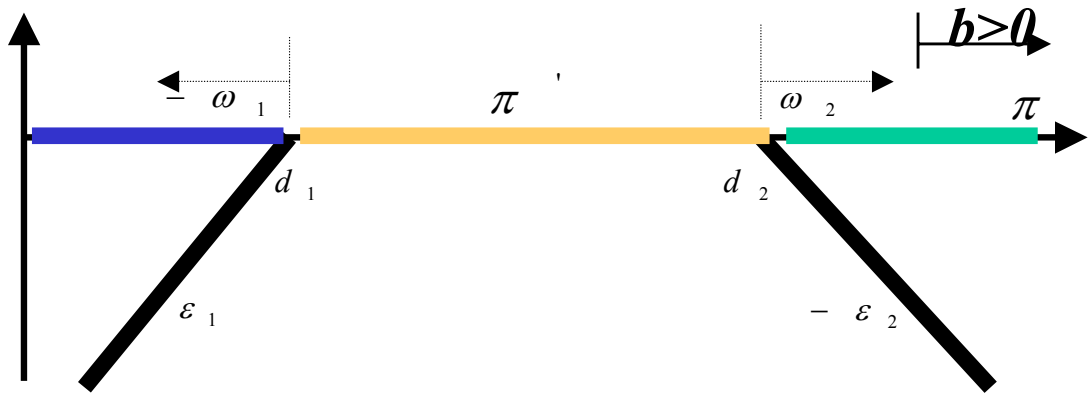
Within a column generation scheme,
problem P is optimal while solving P_ε^d if

$$y_1 = y_2 = 0$$

and

$$\bar{c}_p \geq 0, \text{ for all columns } p \in \Omega$$

Parameter Adjustment, by component

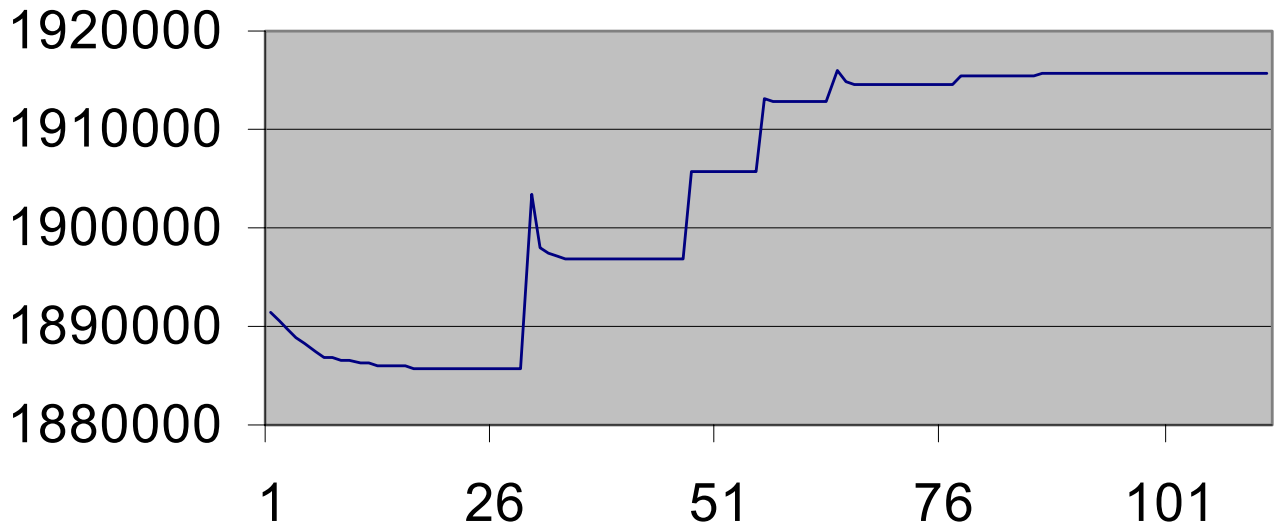


**If the dual value is too small,
re-center and enlarge the interval.**

**If dual value is within the interval,
re-center and reduce the interval.**

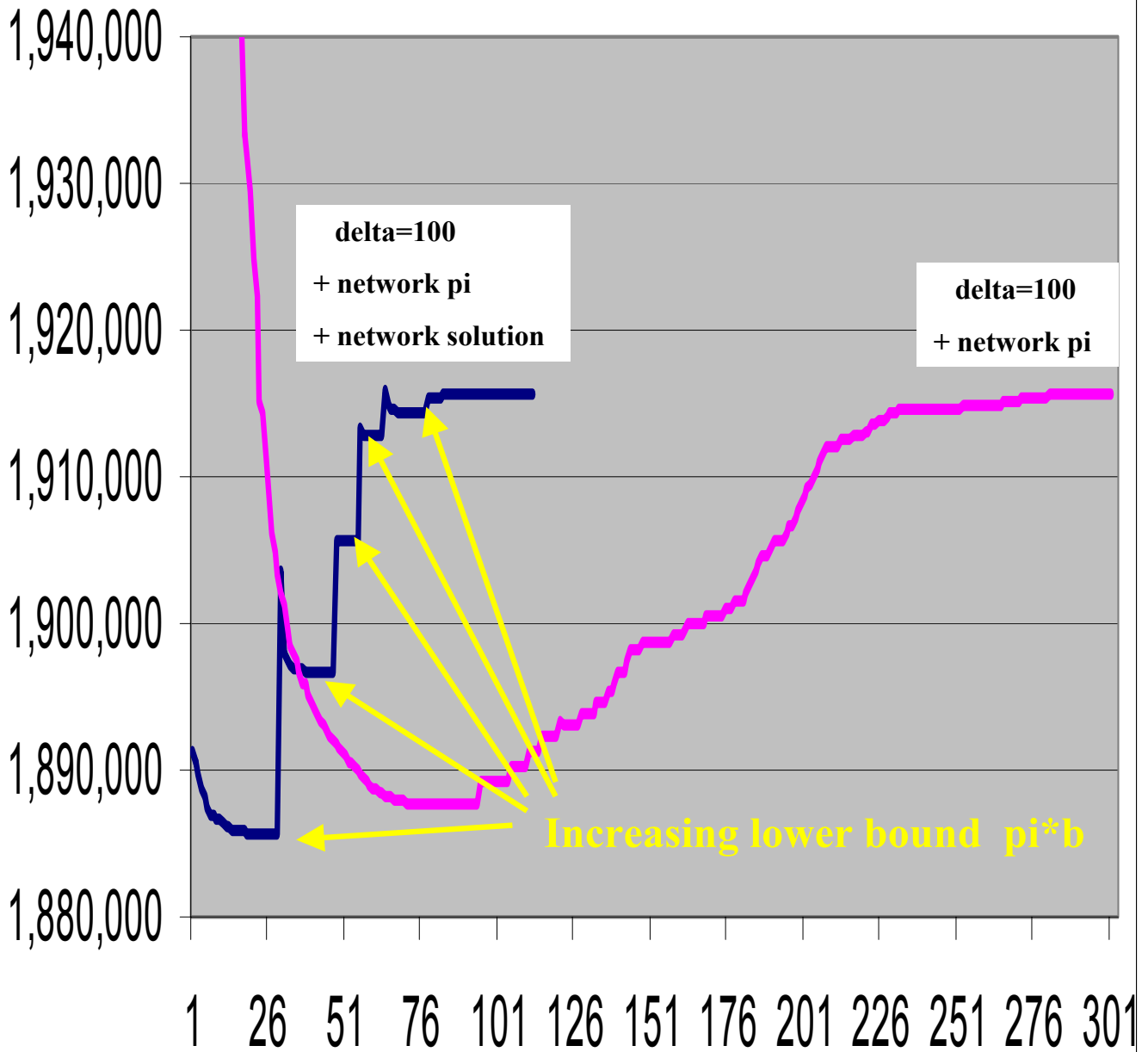
**If the dual value is too large,
re-center and enlarge the interval.**

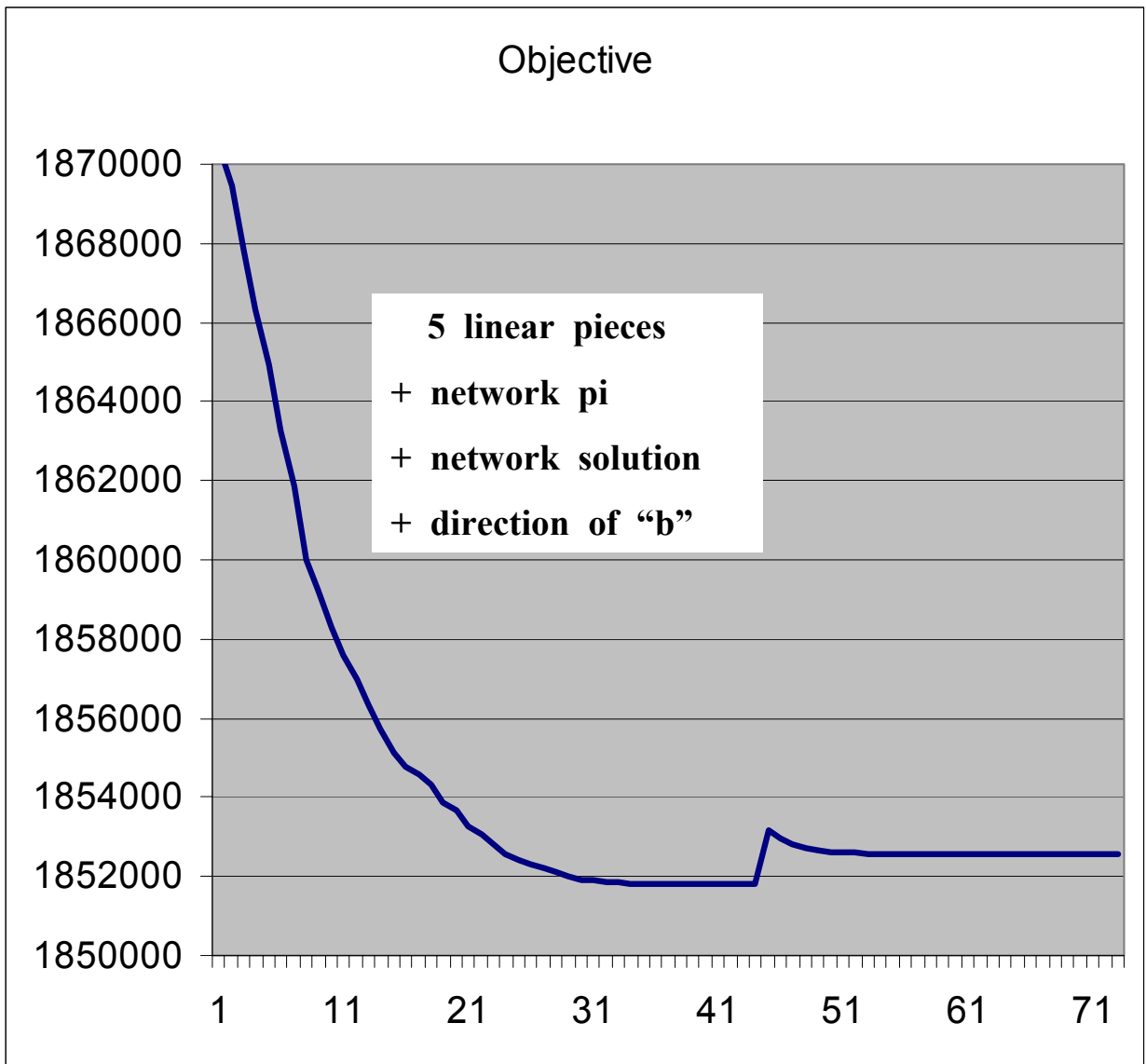
800 Objective Function



<i>Problem</i>	Opt sol	Init sol	<i>cpu tot</i>	cpu mp	cpu sp	# CG iter	# SP cols	# MP itr
<i>R800 (4) standard</i>	1915589.5	800000000	<i>4178.4</i>	3149.2	1029.2	509	37579	926161
delta = 100		2035590.5	<i>835.5</i>	609.1	226.4	119	9368	279155
network pi		2014429.8	<i>1097.1</i>	518.5	578.6	301	10105	324959
network pi + network sol		1891386.0	<i>439.2</i>	216.2	223.0	112	4749	153420
		% reduction	89.5	93.1	78.3	78.0	87.4	83.4
			9.5 times faster					

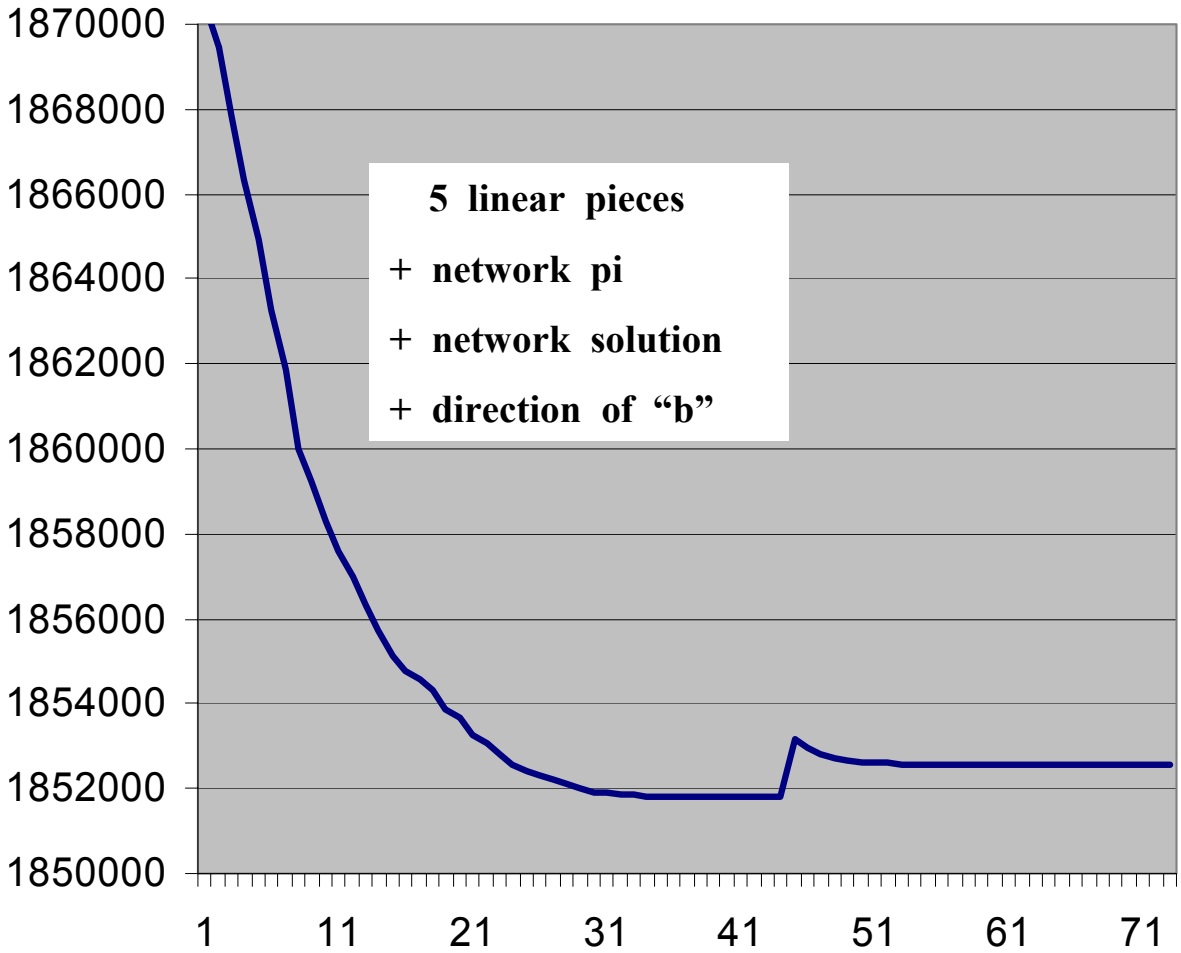
800 Objective Function





<i>Problem</i> P800 (4)	Opt sol	Init sol	<i>cpu tot</i>	cpu mp	cpu sp	# CG iter	# SP cols	# MP it
standard	1852571.5	800000000	3562	886	2676	422	33280	672726
		1870487.8	268	111	157	73	4616	79860
		% reduction	92.5	87.5	94.1	82.7	86.1	88.1
		13.3	times faster					

Objective



Opt sol	Init sol	<i>cpu tot</i>	cpu mp	cpu sp	# CG iter	# SP cols	# MP itr
1852571.5	800000000	<i>3562</i>	886	2676	<i>422</i>	33280	672726
	1870487.8	<i>241</i>	110	131	<i>63</i>		
	% reduction	<i>93.2</i>	<i>87.6</i>	<i>95.1</i>	<i>85.1</i>		
		14.8	times faster				